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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
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EXAMINER	
LE, MIRANDA	
ART UNIT	PAPER NUMBER
2169	
MAIL DATE	DELIVERY MODE
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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

Office Action Summary	Application No.	Applicant(s)	
	10/800,163	DESHMUKH ET AL.	
	Examiner	Art Unit	
	MIRANDA LE	2169	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

1)☒ Responsive to communication(s) filed on *27 August 2008*.

2a)☒ This action is **FINAL**. 2b)☐ This action is non-final.

3)☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

4)☒ Claim(s) *1-18,28,32,34-36 and 39-41* is/are pending in the application.

4a) Of the above claim(s) _____ is/are withdrawn from consideration.

5)☐ Claim(s) _____ is/are allowed.

6)☒ Claim(s) *1-18,28,32,34-36 and 39-41* is/are rejected.

7)☐ Claim(s) _____ is/are objected to.

8)☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

9)☐ The specification is objected to by the Examiner.

10)☐ The drawing(s) filed on _____ is/are: a)☐ accepted or b)☐ objected to by the Examiner.

Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).

Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).

11)☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

12)☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).

a)☐ All b)☐ Some * c)☐ None of:

1.☐ Certified copies of the priority documents have been received.

2.☐ Certified copies of the priority documents have been received in Application No. _____.

3.☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

* See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

1) ☒ Notice of References Cited (PTO-892)
2) ☐ Notice of Draftsperson's Patent Drawing Review (PTO-948)
3) ☒ Information Disclosure Statement(s) (PTO/SB/08) Paper No(s)/Mail Date *08/27/08*.
4) ☐ Interview Summary (PTO-413) Paper No(s)/Mail Date. _____.
5) ☐ Notice of Informal Patent Application
6) ☐ Other: _____.

DETAILED ACTION

This communication is responsive to Amendment, filed 08/27/08.

Claims 1-18, 28, 32, 34-36, 39-41 are pending in this application. This action is made Final.

The objection to claims 32, 34 has been withdrawn in view of the amendment.

Information Disclosure Statement

Applicants' Information Disclosure Statement, filed 08/27/2008, has been received, entered into the record, and considered. See attached form PTO-1449.

Claim Rejections - 35 USC § 102

The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

> A person shall be entitled to a patent unless:
>
> (e) the invention was described in
>
> (1) an application for patent, published under section 122(b), by another filed in the United States before the invention by the applicant for patent or
>
> (2) a patent granted on an application for patent by another filed in the United States before the invention by the applicant for patent, except that an international application filed under the treaty defined in section 351(a) shall have the effects for purposes of this subsection of an application filed in the United States only if the international application designated the United States and was published under Article 21(2) of such treaty in the English language.

Claim 41 is rejected under 35 U.S.C. 102(e) as being anticipated by Horn (US Patent No 7,275,063).

Horn anticipated independent claim 1 by the following:

As per claim 41, Horn teaches a method for storing results of a file walk of a storage server *(i.e. Many of the different instances of information that a user may want to organize are already stored as different types of files in the file system, or otherwise external to the user's client computer or the user's MFS, such as web pages on the World Wide Web or records in an online database. Some information is stored one-for-one: that is, a single file represents a single piece of information (say, a text file). Other information is aggregated into a single file, or is spread across multiple files: for example, email messages are typically stored many to a file due to their small size, and records in a database may be stored across multiple files as well, col. 12, lines 19-37)* comprising:

performing a file walk *(i.e. walking up the outline item's parent tree until there are no more parents, col. 33, lines 3-33; a scanner thread, col. 34, lines 29-65)* of a storage server *(i.e. such as web pages on the World Wide Web or records in an online database, col. 12, lines 19-37)*, wherein performing the file walk includes assigning unique identification numbers *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* to directories *(i.e. Object: means any piece of information stored in a digital format,*

including but not limited to file system entities such as files and folders, col. 4, lines 51-58) of the storage server in a depth first search *(i.e. a depth in the folder hierarchy, col. 34, lines 29-65)* order during the file walk *(i.e. Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array, col. 34, lines 29-65)*;

storing indications *(i.e. a depth in the folder hierarchy; and a flag that determines whether the file system specifier is for a folder or a file, col. 34, lines 29-65; The link metadata, including the UID and UUID, col. 4, lines 39-58; The system, as metadata are created upon selection or creation of collections or containers, "reflects" the reference object in them through tagging additional "path" and "hierarchy" link metadata to the properties metadata that is automatically associated with the reference object and stored in the catalog, col. 8, lines 1-39; Each object has a depth, a numeric value that describes how far down the hierarchy it exists; in particular, how many nodes down the hierarchy tree from the root. Objects at the same depth are known as siblings. The depth determines how far the object is indented to the right in the outline display, col. 32, lines 18-31)* of the directories of the storage server, the indications stored in association with the unique identification numbers *(i.e. Object Store: means a special database that stores and retrieves object data by unique identifier (UID), col. 4, lines 59-60; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* assigned during the file walk *(i.e. Traverse, col. 34, line 29-65)*; and

traversing the stored indications based on the unique identification numbers to determine relationships between the directories of the storage server *(i.e. This thread performs the following functions, in order: 1) Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array, col. 34, lines 29-65)*.

Claim Rejections - 35 USC § 103

The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

> (a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.

This application currently names joint inventors. In considering patentability of the claims under 35 U.S.C. 103(a), the examiner presumes that the subject matter of the various claims was commonly owned at the time any inventions covered therein were made absent any evidence to the contrary. Applicant is advised of the obligation under 37 CFR 1.56 to point out the inventor and invention dates of each claim that was not commonly owned at the time a later invention was made in order for the examiner to consider the applicability of 35 U.S.C. 103(c) and potential 35 U.S.C. 102(e), (f) or (g) prior art under 35 U.S.C. 103(a).

Claims 1-18, 39 are rejected under 35 U.S.C. 103(a) as being unpatentable over Sedlar (US Patent No. 6,922,708), in view of Horn (US Patent No 7,275,063).

As per claim 1, Sedlar teaches a method for creating a file information database *(i.e. Emulating other OS File System Characteristics in a Database, col. 10, line 50 to col. 11, line 7)* comprising:

scanning *(i.e. to scan the table, col. 21, lines 38-49)* a storage server *(i.e. database server 204 stores files that originate from numerous distinct OS file systems, col. 12, lines 21-37)* having a directory structure *(i.e. During the traversal of the hierarchical index, … a child of the directory associated with the directory entry, col. 22, lines 39-47)*;

collecting data regarding the directory structure *(i.e. During the traversal of the hierarchical index, … a child of the directory associated with the directory entry, col. 22, lines 39-47)*;

for each directory of the directory structure, determining whether *(i.e. Once the index entry 508 for the root directory 610 has been located, the DBMS determines whether there are any more filenames in the input pathname (step 802). If there are no more filenames in the input pathname, then control proceeds to step 820 and the FileID stored in index entry 508 is used to look up the root directory entry in the files table 710, col. 9, lines 10-16)* each member of the directory is a file *(i.e. the FileID of directory 616 ("X3"). As shall be described in greater detail, the information contained in the Dir_entry_list field makes accessing information based on pathnames much faster and easier, col. 8, lines 4-17)* or subdirectory *(i.e. the items that have index entries in the hierarchical index 510 are only those directories that are parents to other directories and/or that are currently storing documents. Those items that do not have children (e.g. Example.doc, Access, App1, App2, App3 of FIG. 6) are preferably not included, col. 7, lines 54-64)*;

assigning a first identification (ID) number to a first determined directory *(i.e. index entry has been created for the Documents directory, col. 23, lines 1-8)* and a second ID number to a second determined directory *(i.e. Word directory, col. 23, lines 9-16)* in the directory structure *(i.e. an index entry for the Document directory is added to the hierarchical index, … the Dir_Entry_List is updated to indicate that the new Document directory is a child of the Word directory, col. 22, lines 56-64)*;

examining the determined files *(i.e. Once the index entry 508 for the root directory 610 has been located, the DBMS determines whether there are any more filenames in the input pathname (step 802). If there are no more filenames in the input pathname, then control proceeds to step 820 and the FileID stored in index entry 508 is used to look up the root directory entry in the files table 710, col. 9, lines 10-16)*; and

writing a data structure including the first ID number, the second ID number and relation between the first directory and the second directory *(i.e. an index entry for the Document directory is added to the hierarchical index, … the Dir_Entry_List is updated to indicate that the new Document directory is a child of the Word directory, col. 22, lines 56-64)*.

Sedlar does not explicitly teach:

using first thread to assign a first unique identification (ID) number to a first determined directory and a second unique ID number to a second determined directory in the directory structure according to a depth first search (DFS) order;

wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order;

using second thread to examine the determined files.

Horn teaches:

using first thread to assign a first unique identification (ID) number *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* to a first determined directory and a second unique ID number *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* to a second determined directory in the directory structure according to a depth first search (DFS) order *(i.e. a scanner thread is created with the reference object as a parameter. This thread performs the following functions, in order: 1) Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array. Each entry contains a file system specifier that represents the file; a depth in the folder hierarchy; and a flag that determines whether the file system specifier is for a*

folder or a file. The array is then sorted, deepest objects in the tree first (so that files within a folder are created before the folder is), col. 34, lines 29-65);

wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)*;

using second thread to examine the determined files *(i.e. 2) Annotate: Once this array has been populated, the entries are annotated with metadata that can be efficiently fetched "en masse", such as file and folder comments. 3) Create: The entries are fetched one by one from the array. For each entry, a reference object is created with the entry's information (e.g. the file specifier and any metadata that was previously fetched and added to the catalog). A new array of reference objects is created. 4) Classify: Each object in the array is then classified by examining its metadata and determining in which collections the object belongs, based on the collections' specifications. Every collection that is modified (e.g. that has received a new object through the classification process) is added to yet a third array for notification. 5) Notify: Finally, each collection that participated in the classify step is notified that it has been changed. This typically results in the collection updating dependent property values*

(e.g. count of contained objects), which are then updated in a separate thread, col. 34, lines 29-65).

It would have been obvious to one of ordinary skill of the art having the teaching of Sedlar, Horn at the time the invention was made to modify the system of Sedlar to include the limitations as taught by Horn. One of ordinary skill in the art would be motivated to make this combination in order to assign an internal unique identifier to an object in view of Horn (col. 8, lines 1-39), as doing so would give the added benefit of automatically organizing, indexing and viewing information objects from multiple sources as taught by Horn (col. 6, lines 24-39).

As per claim 10, Sedlar teaches a machine readable medium having stored thereon executable program code which, when executed, causes a machine to perform a method for creating a file information database *(i.e. Emulating other OS File System Characteristics in a Database, col. 10, line 50 to col. 11, line 7)*, the method comprising:

scanning *(i.e. to scan the table, col. 21, lines 38-49)* a storage server *(i.e. database server 204 stores files that originate from numerous distinct OS file systems, col. 12, lines 21-37)* having a directory structure *(i.e. During the traversal of the hierarchical index, … a child of the directory associated with the directory entry, col. 22, lines 39-47)*;

collecting data regarding the directory structure *(i.e. During the traversal of the hierarchical index, … a child of the directory associated with the directory entry, col. 22, lines 39-47)*;

for each directory of the directory structure, determining whether *(i.e. Once the index entry 508 for the root directory 610 has been located, the DBMS determines whether there are any more filenames in the input pathname (step 802). If there are no more filenames in the input pathname, then control proceeds to step 820 and the FileID stored in index entry 508 is used to look up the root directory entry in the files table 710, col. 9, lines 10-16)* each member of the directory is a file *(i.e. the FileID of directory 616 ("X3"). As shall be described in greater detail, the information contained in the Dir_entry_list field makes accessing information based on pathnames much faster and easier, col. 8, lines 4-17)* or subdirectory *(i.e. the items that have index entries in the hierarchical index 510 are only those directories that are parents to other directories and/or that are currently storing documents. Those items that do not have children (e.g. Example.doc, Access, App1, App2, App3 of FIG. 6) are preferably not included, col. 7, lines 54-64)*;

assigning a first identification (ID) number to a first determined directory *(i.e. index entry has been created for the Documents directory, col. 23, lines 1-8)* and a second ID number to a second determined directory *(i.e. Word directory, col. 23, lines 9-16)* in the directory structure *(i.e. an index entry for the Document directory is added to the hierarchical index, … the Dir_Entry_List is updated to indicate that the new Document directory is a child of the Word directory, col. 22, lines 56-64)*;

examining the determined files *(i.e. Once the index entry 508 for the root directory 610 has been located, the DBMS determines whether there are any more filenames in the input pathname (step 802). If there are no more filenames in the input*

pathname, then control proceeds to step 820 and the FileID stored in index entry 508 is used to look up the root directory entry in the files table 710, col. 9, lines 10-16); and

writing a data structure including the first ID number, the second ID number and a relation between the first directory and the second directory *(i.e. an index entry for the Document directory is added to the hierarchical index, … the Dir_Entry_List is updated to indicate that the new Document directory is a child of the Word directory, col. 22, lines 56-64).*

Sedlar does not explicitly teach:

using first thread to assign a first unique identification (ID) number to a first determined directory and a second unique ID number to a second determined directory in the directory structure according to a depth first search (DFS) order;

wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order;

using second thread to examine the determined files.

Horn teaches:

using first thread to assign a first unique identification (ID) number *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* to a first determined directory and a

second unique ID number *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* to a second determined directory in the directory structure according to a depth first search (DFS) order *(i.e. a scanner thread is created with the reference object as a parameter. This thread performs the following functions, in order: 1) Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array. Each entry contains a file system specifier that represents the file; a depth in the folder hierarchy; and a flag that determines whether the file system specifier is for a folder or a file. The array is then sorted, deepest objects in the tree first (so that files within a folder are created before the folder is), col. 34, lines 29-65)*;

wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)*;

using second thread to examine the determined files *(i.e. 2) Annotate: Once this array has been populated, the entries are annotated with metadata that can be efficiently fetched "en masse", such as file and folder comments. 3) Create: The entries are fetched one by one from the array. For each entry, a reference object is created with the entry's information (e.g. the file specifier and any metadata that was previously fetched and added to the catalog). A new array of reference objects is created. 4) Classify: Each object in the array is then classified by examining its metadata and determining in which collections the object belongs, based on the collections' specifications. Every collection that is modified (e.g. that has received a new object through the classification process) is added to yet a third array for notification. 5) Notify: Finally, each collection that participated in the classify step is notified that it has been changed. This typically results in the collection updating dependent property values (e.g. count of contained objects), which are then updated in a separate thread, col. 34, lines 29-65).*

It would have been obvious to one of ordinary skill of the art having the teaching of Sedlar, Horn at the time the invention was made to modify the system of Sedlar to include the limitations as taught by Horn. One of ordinary skill in the art would be motivated to make this combination in order to assign an internal unique identifier to an object in view of Horn (col. 8, lines 1-39), as doing so would give the added benefit of automatically organizing, indexing and viewing information objects from multiple sources as taught by Horn (col. 6, lines 24-39).

As per claim 39, Sedlar teaches a method for creating a file information database *(i.e. Emulating other OS File System Characteristics in a Database, col. 10, line 50 to col. 11, line 7)* comprising:

scanning *(i.e. to scan the table, col. 21, lines 38-49)* a storage server *(i.e. database server 204 stores files that originate from numerous distinct OS file systems, col. 12, lines 21-37)* having a directory structure *(i.e. During the traversal of the hierarchical index, … a child of the directory associated with the directory entry, col. 22, lines 39-47)*;

for each directory of the directory structure, determining whether *(i.e. Once the index entry 508 for the root directory 610 has been located, the DBMS determines whether there are any more filenames in the input pathname (step 802). If there are no more filenames in the input pathname, then control proceeds to step 820 and the FileID stored in index entry 508 is used to look up the root directory entry in the files table 710, col. 9, lines 10-16)* each member of the directory is a file *(i.e. the FileID of directory 616 ("X3"). As shall be described in greater detail, the information contained in the Dir_entry_list field makes accessing information based on pathnames much faster and easier, col. 8, lines 4-17)* or subdirectory *(i.e. the items that have index entries in the hierarchical index 510 are only those directories that are parents to other directories and/or that are currently storing documents. Those items that do not have children (e.g. Example.doc, Access, App1, App2, App3 of FIG. 6) are preferably not included, col. 7, lines 54-64)*;

assigning a first identification (ID) number to a first determined directory *(i.e. index entry has been created for the Documents directory, col. 23, lines 1-8)* and a second ID number to a second determined directory *(i.e. Word directory, col. 23, lines 9-16)* in the directory structure *(i.e. an index entry for the Document directory is added to the hierarchical index, … the Dir_Entry_List is updated to indicate that the new Document directory is a child of the Word directory, col. 22, lines 56-64)*;

examining the determined files *(i.e. Once the index entry 508 for the root directory 610 has been located, the DBMS determines whether there are any more filenames in the input pathname (step 802). If there are no more filenames in the input pathname, then control proceeds to step 820 and the FileID stored in index entry 508 is used to look up the root directory entry in the files table 710, col. 9, lines 10-16)*; and

writing a data structure including the first ID number, the second ID number and relation between the first directory and the second directory *(i.e. an index entry for the Document directory is added to the hierarchical index, … the Dir_Entry_List is updated to indicate that the new Document directory is a child of the Word directory, col. 22, lines 56-64)*.

Sedlar does not explicitly teach:

using first thread to assign a first unique identification (ID) number to a first determined directory and a second unique ID number to a second determined directory in the directory structure according to a depth first search (DFS) order;

wherein the directory numbers are chronologically assigned in numerical order while the directory structure is being traversed in the DFS order;

using second thread to examine the determined files.

Horn teaches:

using first thread to assign a first unique identification (ID) number *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* to a first determined directory and a second unique ID number *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39; B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* to a second determined directory in the directory structure according to a depth first search (DFS) order *(i.e. a scanner thread is created with the reference object as a parameter. This thread performs the following functions, in order: 1) Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array. Each entry contains a file system specifier that represents the file; a depth in the folder hierarchy; and a flag that determines whether the file system specifier is for a*

folder or a file. The array is then sorted, deepest objects in the tree first (so that files within a folder are created before the folder is), col. 34, lines 29-65);

wherein the directory numbers are chronologically assigned in numerical order *(i.e. B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* while the directory structure is being traversed in the DFS order *(i.e. MFS creates an internal representation of the object and stores the representation in the MFS object oriented database (OODB), called the object store, which assigns an internal unique identifier (UID), col. 8, lines 1-39)* ;

using second thread to examine the determined files *(i.e. 2) Annotate: Once this array has been populated, the entries are annotated with metadata that can be efficiently fetched "en masse", such as file and folder comments. 3) Create: The entries are fetched one by one from the array. For each entry, a reference object is created with the entry's information (e.g. the file specifier and any metadata that was previously fetched and added to the catalog). A new array of reference objects is created. 4) Classify: Each object in the array is then classified by examining its metadata and determining in which collections the object belongs, based on the collections' specifications. Every collection that is modified (e.g. that has received a new object through the classification process) is added to yet a third array for notification. 5) Notify: Finally, each collection that participated in the classify step is notified that it has been changed. This typically results in the collection updating dependent property values*

(e.g. count of contained objects), which are then updated in a separate thread, col. 34, lines 29-65).

It would have been obvious to one of ordinary skill of the art having the teaching of Sedlar, Horn at the time the invention was made to modify the system of Sedlar to include the limitations as taught by Horn. One of ordinary skill in the art would be motivated to make this combination in order to assign an internal unique identifier to an object in view of Horn (col. 8, lines 1-39), as doing so would give the added benefit of automatically organizing, indexing and viewing information objects from multiple sources as taught by Horn (col. 6, lines 24-39).

As to claims 2, 11, Horn teaches the agent has a first file system, and the scanning and collecting by using an agent separate from the storage server *(i.e. Many of the different instances of information that a user may want to organize are already stored as different types of files in the file system, or otherwise external to the user's client computer or the user's MFS, such as web pages on the World Wide Web or records in an online database. Some information is stored one-for-one: that is, a single file represents a single piece of information (say, a text file). Other information is aggregated into a single file, or is spread across multiple files: for example, email messages are typically stored many to a file due to their small size, and records in a database may be stored across multiple files as well, col. 12, lines 19-37).*

As to claims 3, 12, Horn teaches the agent has a first file system, and the storage server has a second file system, and wherein the first file system is different from the second file system *(i.e. Many of the different instances of information that a user may want to organize are already stored as different types of files in the file system, or otherwise external to the user's client computer or the user's MFS, such as web pages on the World Wide Web or records in an online database. Some information is stored one-for-one: that is, a single file represents a single piece of information (say, a text file). Other information is aggregated into a single file, or is spread across multiple files: for example, email messages are typically stored many to a file due to their small size, and records in a database may be stored across multiple files as well, col. 12, lines 19-37).*

As to claims 4, 13, Sedlar teaches the relation indicates that the first directory is an immediate child of the second directory *(i.e. an index entry for the Document directory is added to the hierarchical index, … the Dir_Entry_List is updated to indicate that the new Document directory is a child of the Word directory, col. 22, lines 56-64).*

As to claims 5, 14, Horn teaches assigning further comprises assigning the ID number while colleting the data *(i.e. a scanner thread is created with the reference object as a parameter. This thread performs the following functions, in order: 1) Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array. Each entry contains a file system specifier that represents*

the file; a depth in the folder hierarchy; and a flag that determines whether the file system specifier is for a folder or a file. The array is then sorted, deepest objects in the tree first (so that files within a folder are created before the folder is), col. 34, lines 29-65).

As to claims 6, 15, Sedlar teaches writing the data structure *(i.e. an index entry for the Document directory is added to the hierarchical index, … the Dir_Entry_List is updated to indicate that the new Document directory is a child of the Word directory, col. 22, lines 56-64)* further comprises writing the data structure to a database server *(i.e. Database server, Fig. 3)*.

As to claims 7, 16, Sedlar teaches:

receiving a request to determine the parent of the first directory *(i.e. the pathname resolution process for locating a file within an emulated file system begins by locating the index entry 508 of the root directory 610 (step 800), col. 8, line 56 to col. 9, line 9)*; and

referencing the relation between the first directory and the second directory of the data structure to determine the parent of the first directory *(i.e. the pathname resolution process for locating a file within an emulated file system begins by locating the index entry 508 of the root directory 610 (step 800). Because all pathname resolution operations begin by accessing the root directory's index entry 508, data that indicates the location of the index entry for the root directory 610 (index entry 508) may be*

maintained at a convenient location outside of the hierarchical index 510 in order to quickly locate the index entry 508 of the root directory at the start of every search, col. 8, line 56 to col. 9, line 9).

As to claims 8, 17, Sedlar teaches:

receiving a request to determine an immediate child of the second directory *(i.e. Consulting the Dir_entry_list of index entry 512, the system searches for the next filename in the input pathname (steps 804 and 806). In the present example, the filename "Word" follows the filename "Windows" in the input pathname, col. 9, line 56 to col. 10, line 7);*

searching the data structure to find any relation, including the relation between the first directory and the second directory, which indicates that the second directory is a parent in said relation *(i.e. the system searches the Dir_entry_list of index entry 512 for an array entry for "Word", col. 9, line 56 to col. 10, line 7);* and

determine the immediate child of the second directory based on said any relation *(i.e. Since Word directory 616 is just part of the specified path and not the target, files table 710 is not consulted. Instead, the system uses the RowID (Y3) to locate the index entry 514 for Word directory 616 (step 824), col. 9, line 56 to col. 10, line 7).*

As to claims 9, 18, Horn teaches:

receiving a request to determine a set of ID number of every child of a third directory in the directory structure *(i.e. a scanner thread is created with the reference*

object as a parameter. This thread performs the following functions, in order: 1) Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array. Each entry contains a file system specifier that represents the file; a depth in the folder hierarchy; and a flag that determines whether the file system specifier is for a folder or a file. The array is then sorted, deepest objects in the tree first (so that files within a folder are created before the folder is), col. 34, lines 29-65), wherein the third directory is assigned a third ID number *(i.e. B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)*;

determining fourth ID number of a sibling *(i.e. a scanner thread is created with the reference object as a parameter. This thread performs the following functions, in order: 1) Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array. Each entry contains a file system specifier that represents the file; a depth in the folder hierarchy; and a flag that determines whether the file system specifier is for a folder or a file. The array is then sorted, deepest objects in the tree first (so that files within a folder are created before the folder is), col. 34, lines 29-65)* of the third directory *(i.e. B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52)* ; and

determining the set of ID number between the third ID number *(i.e. a scanner thread is created with the reference object as a parameter. This thread performs the following functions, in order: 1) Traverse: A procedure recursively descends the folder hierarchy, creating data entries that are stored in an array. Each entry contains a file system specifier that represents the file; a depth in the folder hierarchy; and a flag that determines whether the file system specifier is for a folder or a file. The array is then sorted, deepest objects in the tree first (so that files within a folder are created before the folder is), col. 34, lines 29-65)* and the fourth ID number *(i.e. B-Tree: means a data structure by which information may be stored efficiently on disk, with a minimum of disk accesses to fetch a particular piece of information using an ordered key such as a numeric identifier or a sortable string of text, col. 3, lines 48-52).*

Claims 28, 34-36 are rejected under 35 U.S.C. 103(a) as being unpatentable over Chen et al. (US Patent No. 6,625,624), in view of Tan et al. (US Patent No. 6,356,902).

As per claim 28, Chen teaches a method for creating a logical tree comprising:

using a directory walking thread *(i.e. walking thread, See Fig. 2)* to examine a first directory *(i.e. a root URL, col. 7, lines 20-32)* from a top of a directory queue *(i.e. As described above, the walking can be defined by a root URL and parameterized by (a) the depth of walking through hyper-references under HTML pages, (b) with or without image files embedded in HTML pages, and (c) walking through pages on the local web site or on all web sites, col. 7, lines 20-32)* and determine a set of children of the directory *(i.e. archive pages, col. 3, lines 54-60; The iagent class 2035 of the agent*

thread 2030 connects to a remote web server 70 or proxy to request a Web page. iagent 2035 can cache the page and return the page to iserver 2025 or to ihtwalk 2045, a facility to walk the html tree structure to collect and archive pages. The ihtwalk class facility is further described below in Section 4 in the description of the walking facilities, col. 3, lines 54-60);

examining a set of children *(i.e. HTML pages, col. 7, lines 20-32)* of the first directory to determine a first subset of files and a second subset of directories *(i.e. walking through pages on the local web site or on all web sites, col. 7, lines 20-32)*;

placing the first subset of files in a file queue for examination by a file thread *(i.e. the walking facility may be used to visit the sets of web pages that are designated to be packed, col. 8, lines 3-23)*; and

placing the second subset on the top of the directory queue *(i.e. The system then walks through the set of web pages rooted by the designated URL and packs them into a package using the designated name col. 8, lines 3-23).*

Chen does not explicitly teach:

assigning a depth first search (DFS) ID to the first directory, wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order.

Tan teaches:

assigning a depth first search (DFS) ID to the first directory *(i.e. means of traversing in depth-first search in a depth coupled graph map without page fault using an efficient traversing algorithm, col. 2, lines 46-59)*, wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order *(i.e. For each*

level of the tree structure, the graph nodes are stored in the Parent-stack, the adjacent level of tree structure is stored in the child-stack. Assignment of the link pointer of the graph node in the parent-stack is stated in the conditions shown in step 2.9, col. 5, line 7 to col. 6, line 3).

It would have been obvious to one of ordinary skill of the art having the teaching of Chen and Tan at the time the invention was made to modify the system of Chen to include the limitations as taught by Tan. One of ordinary skill in the art would be motivated to make this combination in order to traverse in depth-first search in a depth couple graph map in view of Tan (col. 2, lines 46-59), as doing so would give the added benefit of enhancing data retrieval and storage in a multilevel tree structure by significantly reducing the required amount of time and memory as taught by Tan (col. 2, lines 40-45).

As per claim 34, Chen teaches the method of claim 28, wherein the directory walking thread *(i.e. walking thread, See Fig. 2)* is hosted by an agent *(i.e. Agent Thread, See Fig. 2)* that is separate from the storage server *(See Fig. 1).*

As per claim 35, Chen teaches the method of claim 34, further comprising using an MMA *(i.e. IProxy main thread 2010, Fig. 2)* to control the agent *(See Figs. 1, 2).*

As per claim 36, Chen teaches the method of claim 34, wherein the directories are hosted by a filer *(i.e. Agent Thread, See Fig. 2).*

Claim 40 is rejected under 35 U.S.C. 103(a) as being unpatentable over Sedlar (US Patent No. 6,922,708), in view of Horn (US Patent No 7,275,063), and further in view of Gasser (US Patent No. 6,636,250).

As per claim 40, Sedlar, Horn do not explicitly teach the method of claim 1, wherein a top level directory of the directory structure is assigned an ID of "0" (zero).

Gasser teaches this limitation in Fig. 9.

It would have been obvious to one of ordinary skill of the art having the teaching of Sedlar, Horn, Gasser at the time the invention was made to modify the system of Sedlar, Horn to include the limitations as taught by Gasser. One of ordinary skill in the art would be motivated to make this combination in order define in a first level of the tree serve as group arrangement descriptors in view of Gasser (col. 5, lines 30-57), as doing so would give the added benefit of allowing a user to overlay multiple relationships on top of one another so that the user can determine and view the relationships at the same time in the same view of the graphical user interface as taught by Gasser (col. 2, line 48 to col. 8, line 9).

Claim 32 is rejected under 35 U.S.C. 103(a) as being unpatentable over Chen et al. (US Patent No. 6,625,624), in view of Tan et al. (US Patent No. 6,356,902), and further in view of Sedlar (US Patent No. 6,922,708).

As per claim 32, Chen and Tan do not explicitly teach the method of claim 28, wherein examining the file queue further comprises recording information about a first file taken from the file queue.

Sedlar teaches wherein examining the file queue further comprises recording information about a first file taken from the file queue *(i.e. According to one such embodiment, an event server executing external to a database server is registered as a subscriber to a queue managed by the database server. The queue to which the event server subscribes shall be referred to herein as the file event queue. Entities that are interested in particular file system events register their interest with the event server. The event server communicates with the database server through the database API, and with the interested entities through the protocols supported by those entities, col. 28, lines 39-47).*

It would have been obvious to one of ordinary skill of the art having the teaching of Chen, Tan, Sedlar at the time the invention was made to modify the system of Chen, Tan to include the limitations as taught by Sedlar. One of ordinary skill in the art would be motivated to make this combination in order to communicate with the database server in view of Sedlar (col. 28, lines 39-47), as doing so would give the added benefit of the process of formulating and submitting queries to a database server is less intuitive than merely traversing a hierarchy of directories, and is beyond the technical comfort level of many computer users as taught by Sedlar (col. 3, lines 4-12).

Response to Arguments

Applicant's arguments filed 08/27/2008 have been fully considered but they are not persuasive.

1. Claims 1, 10, 39 have been amended to specify the term "unique identifier"; thus, Applicant's argument with respect to " Sedlar, Tan, and Chen do not teach "using a first

thread to assign a first unique identification (ID) number to a first determined directory and a second unique ID number to a second determined directory; using a second thread to examine the determined files; wherein the directory numbers are assigned while the directory structure is being traversed" has been considered but are moot in new ground(s) of rejections.

2. Regarding claim 28:

a. In response to Applicant's argument with respect to "Chen and Tan do not teach assigning a depth first search ID to the first directory, wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order" or placing the first subset of files in a file queue for examination by a file thread, the Examiner respectfully disagrees.

Chen teaches a depth first search as shown in Fig. 4 (i.e. depth = 1 lever).

Chen teaches a depth first search ID as <http://interactive.wsj.com/pages/techman> and level 1 (See Fig. 4).

Chen teaches the directory numbers are assigned before the directory structure is traversed (See Fig. 4).

Chen does not explicitly teach assigning a depth first search ID to the first directory, wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order.

Tan teaches the step of assigning a depth first search ID to the first directory, wherein the directory numbers are assigned while the directory structure is being traversed in the DFS order as in Fig. 2A.

It should be noted that a depth first search ID to the first directory limitation equates to LEV = 1 of Tan in Fig. 2A.

b. In response to Applicant's argument with respect to "Chen and Tan do not teach placing the first subset of files in a file queue for examination by a file thread", the Examiner respectfully traverses.

Chen teaches a walking thread in Fig. 2 or waking action in col. 5, lines 60 to col. 6, line 4.

Chen teaches a file thread corresponding to agent thread in Fig. 2.

Chen teaches the step of placing the first subset of files in a file queue for examination by a file thread as *one or more of a list of functions can be invoked one by one to perform tasks on a cache of the page (i.e. One useful facility is a basic function supporting a mechanism to walk through document page hierarchies. This is similar to what "find" and "tw" (file tree walk) do on Unix file systems. The walking action is specified by a root URL where the action starts, a specification of how many levels the action will visit, and certain additional properties such as whether or not image files should be included. For each visited page, one or more of a list of functions can be invoked one by one to perform tasks on a cache of the page. Examples of such*

functions include functions for archiving the web pages, searching for keywords, and creating index tables, col. 5, line 60 to col. 6, line 4, Chen).

Based on the foregoing, it is submitted that all claims are not patentably distinct over the cited art of record.

Conclusion

9. Applicant's amendment necessitated the new ground(s) of rejection presented in this Office action. Accordingly, **THIS ACTION IS MADE FINAL.** See MPEP § 706.07(a). Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire THREE MONTHS from the mailing date of this action. In the event a first reply is filed within TWO MONTHS of the mailing date of this final action and the advisory action is not mailed until after the end of the THREE-MONTH shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than SIX MONTHS from the date of this final action.

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Miranda Le whose telephone number is (571) 272-4112. The examiner can normally be reached on Monday through Friday from 10:00 AM to 6:00 PM.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, James K. Trujillo, can be reached on (571) 272-3677. The fax number to this Art Unit is (571)-273-8300.

Any inquiry of a general nature or relating to the status of this application should be directed to the Group receptionist whose telephone number is (571) 272-2100.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free).

/Miranda Le/
Primary Examiner, Art Unit 2169